

Parallelizing NEC's Equation Solver Algorithm with OpenMP

Mario Trangoni¹ Victor Rosales²

Argentina Software Design Center (Intel)

¹mario.trangoni@intel.com

²victor.h.rosales@intel.com



HPCLatAm 2012
Buenos Aires, Argentina
July 24th, 2012

Agenda

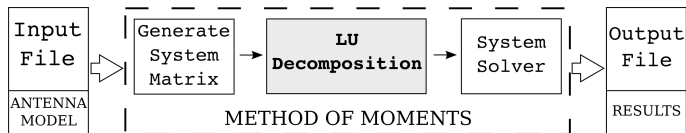
- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation
- 4 Parallelization Proposal
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

Agenda

- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation
- 4 Parallelization Proposal
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

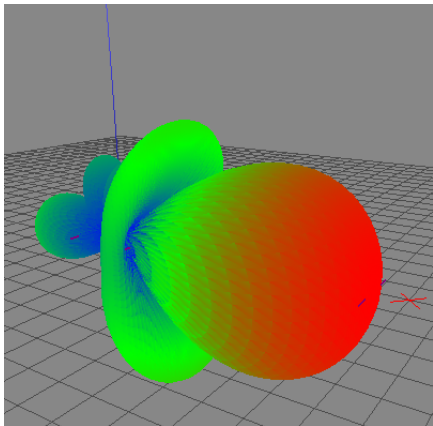
Introduction I

- ▶ NEC was written in FORTRAN in the 1970s and translated from FORTRAN to C in 2003 renamed as *nec2c*.
- ▶ It uses the method of moments solution of the electric field integral equation for thin wires and the magnetic field integral equation for closed, conducting surfaces.
- ▶ Models are defined as elements of wire or similar as an input text file. These models are then input into the NEC application.

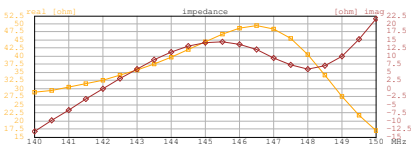
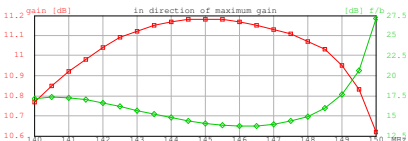
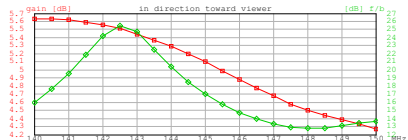


Introduction II

- ▶ Results can then be input into subsequent 'helper' applications for visual viewing and the generation of other graphical representations.



qantenna 0.2



xnecview 1.35

Agenda

- 1 Introduction
- 2 **Discovering Bottlenecks**
- 3 Original Implementation
- 4 Parallelization Proposal
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

Discovering Bottlenecks

- ▶ Intel[®] VTune Amplifier was used to detect the hot spots.
- ▶ nec2c spends most of the execution time in *factr* (LU factorization)
- ▶ *mulxc3* does complex numbers multiplication and is invoked mostly by *factr*.

Hotspots - Hotspots

Analysis Target | Analysis Type | Collection Log | Summary | Bottom-up | Top-down Tree

Grouping: Function / Call Stack

Function / Call Stack	CPU Time	Module	Function (Full)
factr	40.977s	nec2c	factr
factrs ← main ← _start	40.977s	nec2c	factr
_mulxc3	28.325s	nec2c	_mulxc3
factr ← factrs ← main ← _start	27.545s	nec2c	_mulxc3
ffld ← rdpat ← main ← _start	0.280s	nec2c	_mulxc3
eksc ← efld ← cmww ← cmset ← main ← _start	0.229s	nec2c	_mulxc3
gx ← eksc ← efld ← cmww ← cmset ← main	0.171s	nec2c	_mulxc3
solve ← solves ← netwk ← main ← _start	0.080s	nec2c	_mulxc3
efld ← cmww ← cmset ← main ← _start	0.010s	nec2c	_mulxc3
rdpat ← main ← _start	0.010s	nec2c	_mulxc3
ffld	6.464s	nec2c	ffld
gf	2.467s	nec2c	gf
gx	1.228s	nec2c	gx
efld	1.190s	nec2c	efld
eksc	1.118s		
cmww	0.769s		
intx	0.682s		
test	0.250s		

Elapsed Time: 84.094s

Total Thread Count: 1
CPU Time: 84.090s
Paused Time: 0s

Agenda

- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation**
- 4 Parallelization Proposal
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

Original Implementation

- ▶ Works one row at a time from top to bottom

Row 0

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Division: $a_{12} = \frac{a_{12}}{a_{11}}$; $a_{13} = \frac{a_{13}}{a_{11}}$

Row 1

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Subtraction: $a_{22} = a_{22} - a_{12} \cdot a_{21}$;
 $a_{23} = a_{23} - a_{13} \cdot a_{21}$

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Division: $a_{23} = \frac{a_{23}}{a_{22}}$

Row 2

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Subtraction: $a_{32} = a_{32} - a_{12} \cdot a_{31}$;
 $a_{33} = a_{33} - a_{13} \cdot a_{31}$

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Subtraction: $a_{33} = a_{33} - a_{23} \cdot a_{32}$

Agenda

- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation
- 4 Parallelization Proposal**
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

Parallelization Proposal

Matrix (3×3):

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Division: $a_{12} = \frac{a_{12}}{a_{11}}$; $a_{13} = \frac{a_{13}}{a_{11}}$

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Subtraction: $a_{22} = a_{22} - a_{12} \cdot a_{21}$;
 $a_{23} = a_{23} - a_{13} \cdot a_{21}$

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Subtraction: $a_{32} = a_{32} - a_{12} \cdot a_{31}$;
 $a_{33} = a_{33} - a_{13} \cdot a_{31}$

Matrix (2×2):

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Division: $a_{23} = \frac{a_{23}}{a_{22}}$

$$\begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

Subtraction: $a_{33} = a_{33} - a_{23} \cdot a_{32}$

- ▶ First step operates on the full matrix and the subsequent steps will successively reduce the size of the operated matrix.

Agenda

- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation
- 4 Parallelization Proposal
- 5 Implementation**
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

Implementation I

Maximum Search

- ▶ An implementation with *Maximum Search*, *Row Exchange and Division by Row's Diagonal Element* and *Linear Combinations* subsections was executed and compared against an implementation including only the last two subsections.

Number of Threads	Implementation 1 (sec.)	Implementation 2 (sec.)	Difference (Imp.1 – Imp.2)
2	906.71	870.15	4.20 %
4	531.62	516.49	2.93 %
8	344.77	343.93	0,24 %

Comparison of performance over a matrix of 5000 × 5000 elements.

- ▶ *Maximum Search* parallelization not implemented.

Implementation II

Row Exchange and Division by Rows Diagonal Element

- ▶ Exchange of rows only if the maximum is not located on the diagonal.
- ▶ Avoid division by zero and reduce the precision error.
- ▶ Division by the diagonal element or first element of the row.

```
# pragma omp parallel shared (a,...) private(i,...)
  { /*Begin parallel zone*/
if ((ip[r]!=r+1) && (r!=n-1)) {
    i=ip[r]-1; /*if max not at diagonal, do interchange */
# pragma omp for
    for (j=r; j<n; j++) { /*Row's interchange */
        arj=a[i+j*ndim];
        a[i+j*ndim]=a[r+j*ndim];
        a[r+j*ndim]=arj;
    }
}

# pragma omp for
    for (i=r+1; i<n; i++) { /*n-1 iterations */
        a[i+rndim]=a[i+rndim]/a[r+rndim];
    }
}
```

Implementation III

Linear Combinations

- ▶ Simply parallelized with an OpenMP pragma.
- ▶ Parallelization's maximum gain zone.

```
# pragma omp for
  for (j=r+1; j<n; j++) {
    /*n-1 iterations - rows */
    for (i=r+1; i<n; i++) {
      /*n-1 iterations - elements */
      a[i+j*ndim]-=(a[i+r*ndim]*a[r+j*ndim]);
    }
  }
} /*Finish Parallel zone*/
```

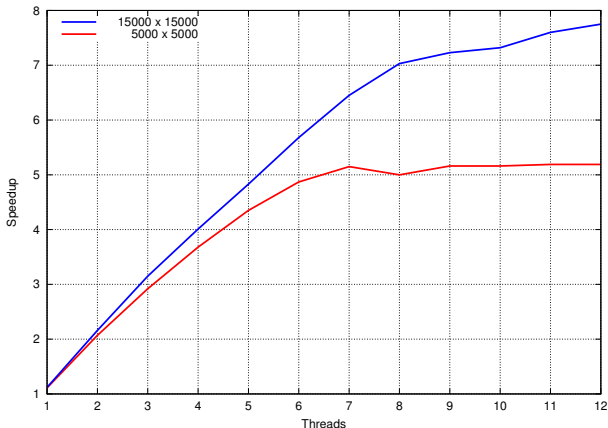
Agenda

- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation
- 4 Parallelization Proposal
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

Results

Scalability over the number of OpenMP threads

- ▶ GCC 4.4.4 - Flags: `-O2 -march=core2 -mtune=generic -fopenmp`
- ▶ Ran on: Intel Westmere-EP (2 x Intel Xeon X5670 (6 cores) @ 2.93 Ghz / 12 MB shared Last-Level Cache / 12Gb RAM DDR3-1066 (8533,33 MB/seg))
- ▶ OS: CentOS release 5.4 (Final)



Agenda

- 1 Introduction
- 2 Discovering Bottlenecks
- 3 Original Implementation
- 4 Parallelization Proposal
- 5 Implementation
 - Maximum Search
 - Row Exchange and Division by Rows Diagonal Element
 - Linear Combinations
- 6 Results
- 7 References

References

Roger F. Harrington, *Field Computation by Moment Methods*, IEEE Press Series on Electromagnetic Wave Theory, Wiley-IEEE, ISBN 0-78-031014-4, 1993.

F. Gisin, *Using the method of moment NEC code to solve EMC problems*, IEEE International Symposium on Electromagnetic Compatibility, Dallas, TX, USA, 9-13 August 1993.

H. Sormann, *Numerische Methoden in der Physik*, Institut für Theoretische Physik - Computational Physics, TU Graz, Graz, Austria, 2011.

G. J. Burke, A. J. Poggio, *Numerical Electromagnetics Code (NEC-2)*, Public Domain Code, 1980.

Thanks!